

A Knowledge-Based Analysis of Intersection Protocols

DISC 2024

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 - Going through intersections should become much more efficient, without the need for traffic lights.
- Fault-tolerance is desirable. Solutions need to handle:
 - communication failures of driverless cars
 - human drivers without wireless communication
- There are many types of intersections:
 - The model should capture all of them, and solutions need to apply broadly.

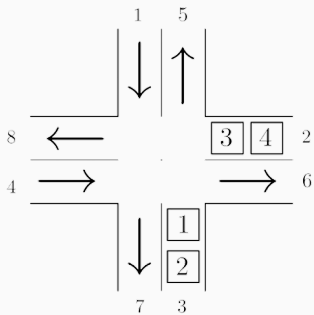
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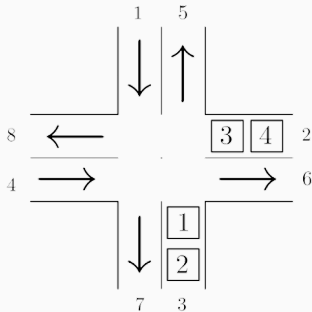
To the best of our knowledge, prior approaches mainly focused on:

- Specific intersection scenarios [RBS21, SSP17].
- Leader-election protocols without communication failures [FVP⁺13, FFCa⁺10].
- Simulation of scenarios [FFCa⁺10, RBS21].

Intersection model

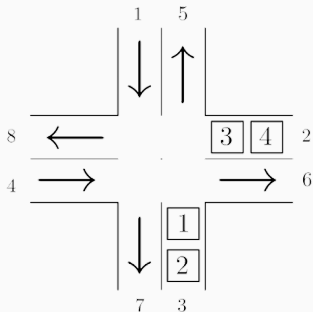


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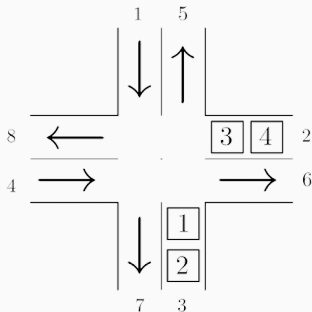
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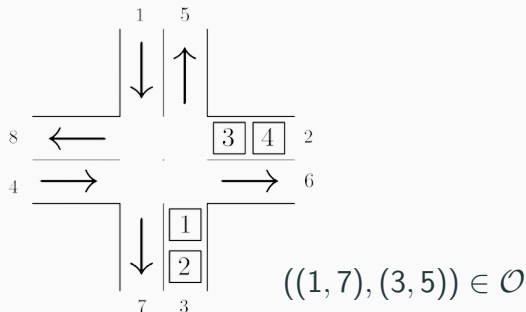
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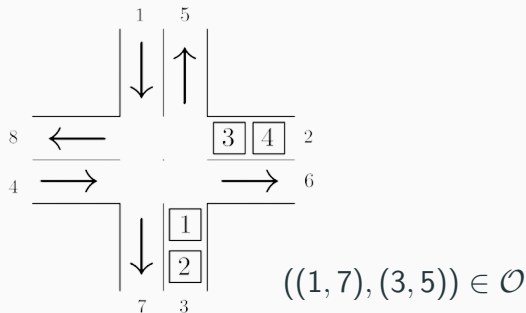
- Vehicles are numbered $Ag = \{1, 2, \dots, \dots\}$ (potentially ∞)
- Incoming lanes are \mathcal{L}_{in} , outgoing lanes are \mathcal{L}_{out} .
- Each vehicle when in front makes a move (ℓ_{in}, ℓ_{out}) .

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- A relation \mathcal{O} determines which moves are *compatible*.

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- The arrival schedule of vehicles and their planned moves are adversarially chosen.

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Our goal is to find protocols that satisfy these properties:

- **Validity:** a vehicle executes **go** only when it is in front
- **Safety:** if two vehicles execute **go**, their moves are compatible
- **Liveness:** if a vehicle is in front, then at some point in the future it executes **go**

*Another way to think about this problem is as a generalization of distributed mutual exclusion

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- we define optimality criteria for intersection protocols:
→ intuitively, vehicles should never wait *unnecessarily*
- we give constructions that result in optimal protocols:
→ via a “global” to “local” reduction

Reasoning about knowledge

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- At a point (r, m) , we can interpret formulas about the state of the system.
- Key definition: $\mathcal{I}, (r, m) \models K_i\phi$ if $\mathcal{I}, (r', m') \models \phi$ at all points (r', m') such that i has the same local state in (r, m) and (r', m')

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Some useful formulas:

- $\mathcal{I}, (r, m) \models \textit{going}_i$ if i executes go at (r, m)
- $\mathcal{I}, (r, m) \models \textit{front}_i$ if i is in front its lane at (r, m)

Optimality criteria

Intuitive criteria:

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Domination-based optimality criteria (from knowledge-based BA literature):

- A protocol is optimal if there is no other protocol that strictly dominates it:
 - P dominates P' if agents in P under the same adversary (i.e. under the same message failures and vehicle arrival schedule) always go through the intersection earlier or at the same time as in P' .

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Eliminating unnecessary waiting is not always possible but lexicographical optimality can be achieved under weaker conditions.

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- The converse is also true (Proposition 4 and 5) if information exchange is “sufficiently rich” and the decision rule only depends on agents in front of each queue.

Intersection policies

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Proposition 7: For every protocol P satisfying Validity and Safety there exists an intersection policy σ such that P implements \mathbf{P}^σ .

Lexicographically optimal protocols

A lexicographically optimal protocol **P**:

if $K_i(\text{front}_i \wedge (i\text{'s move is in } \sigma \vee V_i))$ **then** go
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Intuitively, **P** allows the following moves:

- all moves permitted by σ
- other moves not in σ in cyclic priority order if each agent knows that its move is compatible with the moves of all agents of higher priority (including agents permitted to go by σ).

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*For example, if the precedence cycle in a given point (r, m) is from $k = m \bmod |\mathcal{L}_{in}|$ to $k + 1 \bmod |\mathcal{L}_{in}|, k + 2 \bmod |\mathcal{L}_{in}|, \dots$, and $\sigma = \emptyset$ we satisfy these assumptions.

Future directions

- Tolerating stronger adversaries
- Evaluating implementations in other contexts
- Considering strategic behavior

Thank you!



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